

(47)

$$A \cap B = \{x \in U \mid (x \in A) \wedge (x \in B)\}$$

$$A \cup B = \{x \in U \mid (x \in A) \vee (x \in B)\}$$

$$\overline{A} = \{x \in U \mid \neg(x \in A)\}$$

↗
equiv to writing
" $x \notin A$ "

- we'll explore connections between connectives and set operations more later.

Implication

Given statements

P, Q the statement $P \Rightarrow Q$ is read "if P , then Q " or " P implies Q "

$P \Rightarrow Q$ is true iff

whenever P is true, Q is also true

P	Q	$P \Rightarrow Q$
T	T	T
T	F	F
F	T	T
F	F	T

Notice: - $P \Rightarrow Q$ is always true when P is false (often a confusing point)
- $P \Rightarrow Q$ only false when P is true and Q is false

- statements of the form $P \Rightarrow Q$
are called "conditional statements"

Ex's ① " $(1+1=2) \Rightarrow (1+1+1=3)$ " is true
 $\begin{array}{ccc} \uparrow & & \uparrow \\ \text{T} & & \text{T} \end{array}$

② " $(1+1=2) \Rightarrow (1+1+1=4)$ " is false
 $\begin{array}{ccc} \uparrow & & \uparrow \\ \text{T} & & \text{F} \end{array}$

③ " $(1+1=2) \Rightarrow (\sqrt{2} \notin \mathbb{N})$ " is true

even though P, Q in this example
are apparently unrelated statements.

④ "My name is Sally $\Rightarrow$ My
name begins with S" is true

$\hookrightarrow$ both premise P and conclusion
 Q are false, therefore (by def'n)

$P \Rightarrow Q$ is true ("false $\Rightarrow$ false" is (T))

⑤ "Tomorrow is Sunday $\Rightarrow$ My
name is Garrett" is also true

("false $\Rightarrow$ true" is (T))

⑥ $(\exists x \in \mathbb{R})(x^2 = -1) \Rightarrow (1+1=3)$ (49)

- is true: automatically since premise,
- is false, though unrelated to conclusion

⑦ can also use $\Rightarrow$ in var. prop'n

e.g. $x \geq 2 \Rightarrow x^2 \geq 4$

is a var. prop'n and

$$(\forall x \in \mathbb{R})(x \geq 2 \Rightarrow x^2 \geq 4)$$

is a true statement, since
for every $x \in \mathbb{R}$

either $x \geq 2$ is true,
 in which case $x^2 \geq 4$ is also true
~~and~~ and "true $\Rightarrow$ true" is true
or ~~and~~ $x \geq 2$ is false,
 in which case $x \geq 2 \Rightarrow x^2 \geq 4$
 is true automatically

i.e. for every $x \in \mathbb{R}$, " $(x \geq 2) \Rightarrow (x^2 \geq 4)$ "
 is (T), i.e. $(\forall x \in \mathbb{R})(x \geq 2 \Rightarrow x^2 \geq 4)$
 is (T)

⑧ OTOIT: $(\forall x \in \mathbb{R}) (x^2 \geq 4 \Rightarrow x \geq 2)$ (50)
 is false since there is a real number x (e.g. $x = -3$) s.t. $x^2 \geq 4$ is (T)
 but $x \geq 2$ is (F).

Equivalence Given statements P, Q
 the statement $P \Leftrightarrow Q$
 (read: "P if and only if Q")
 is true iff P, Q have the same truth values

P	Q	$P \Leftrightarrow Q$
T	T	T
T	F	F
F	T	F
F	F	T

Ex's ① $(1+1=2) \Leftrightarrow (1+1+1=3)$
 ② $(1+1=3) \Leftrightarrow (1+1+1=4)$ is (T)
 ③ $(\forall x \in \mathbb{N}) (x > 0) \Leftrightarrow (1+1=2)$ is (T)
 ↗ ↘
 P, Q need not be
 "related"

④ $(1+1=2) \Leftrightarrow (2+2=5)$ is (F)

can also use $\Leftrightarrow$ in var. prop'n, e.g.

$$(x \geq 0) \Leftrightarrow (\exists y \in \mathbb{R}) (y^2 = x)$$

is a legit var. prop'n

and $(\forall x \in \mathbb{R}) [(x \geq 0) \Leftrightarrow (\exists y \in \mathbb{R}) (y^2 = x)]$ (51)
is a true statement, since:

for every $x \in \mathbb{R}$ the ~~the~~ statements
" $x \geq 0$ " and " $(\exists y \in \mathbb{R}) (y^2 = x)$ " are
either both true or both false.

Def'n Statements P, Q are said
to be logically equivalent iff they
have the same truth value, i.e.
iff $P \Leftrightarrow Q$ is true.

- e.g. $1+1=2$ and $1+1+1=3$ are logically equiv.

- will be most interested in
logically equivalent forms for
connected (esp. negated) and
quantified statements.

Negating Quantified statements

- Sp's $P(x)$ is a variable prop'n
and S is a set.

- Consider the negated statements (52)

$$\textcircled{1} \quad \neg (\forall x \in S) P(x)$$

$$\textcircled{2} \quad \neg (\exists x \in S) P(x)$$

Observe $\textcircled{1}$ is true iff there is an $x \in S$ s.t. $P(x)$ is false, i.e. iff $(\exists x \in S) \neg P(x)$ is true

$\textcircled{2}$ is true iff for all $x \in S$ we have $P(x)$ is false, i.e. $(\forall x \in S) \neg P(x)$ is true.

This shows

$$\neg (\forall x \in S) P(x) \Leftrightarrow (\exists x \in S) \neg P(x)$$

is always true (regardless of $P(x)$)

i.e. that $\neg (\forall x \in S) P(x)$ and $(\exists x \in S) \neg P(x)$ are logically equivalent.

Likewise

$$\neg (\exists x \in S) P(x) \text{ and } (\forall x \in S) \neg P(x)$$

are logically equiv.

— these equivalences often useful when trying to prove quantified statements by contradiction.

Ex 5 ① $\neg(\forall x \in \mathbb{R})(x \in \mathbb{N})$

"not all reals are naturals"

is equiv. to

$$(\exists x \in \mathbb{R}) \neg(x \in \mathbb{N})$$

"there is a real which is not a natural"

both are (T)

(note: we'll often write $\neg(x \in \mathbb{N})$

as $x \notin \mathbb{N}$, $\neg(x = y)$ as $x \neq y$, etc.)

② $\neg(\exists x \in \mathbb{R})(x+1=0)$

is equiv to

$$(\forall x \in \mathbb{R})(x+1 \neq 0)$$

(both (F))

③ For multiple quantifiers, iterate the process:

$$\neg(\forall x \in \mathbb{R})(\exists y \in \mathbb{R})(xy=1) \rightarrow$$

"not every real has a mult. inverse"

equiv. to: $(\exists x \in \mathbb{R}) \neg(\exists y \in \mathbb{R})(xy=1)$

$$\text{also: } (\exists x \in \mathbb{R})(\forall y \in \mathbb{R})(xy \neq 1)$$

(both are (T) :

C has no multiplicative
inverse)

→ (54)
"there is a
red w/o a
multiplicative
inverse"