

Def'n A prime factorization of a ⁽⁹⁶⁾
given $n \in \mathbb{N}$ ^{$n > 1$} is a way of writing
 n as a product of primes.

e.g. $6 = 2 \cdot 3$ is a p.f. of 6

$2 = 2$ is a p.f. of 2

$100 = 2 \cdot 2 \cdot 5 \cdot 5$ is a p.f. of 100

Thm For every $n \in \mathbb{N}$, $n > 1$ (i.e. $n \in \{2, 3, 4, \dots\}$),
 n has a prime factorization.

Pf: Let $P(n)$ be the prop'n

" n has a prime factorization"

(BC) $P(2)$ holds since 2 has a p.f.

(IH) Fix $n \in \{2, 3, 4, \dots\}$ and assume
that $\forall k \in \{2, 3, \dots, n\}$ k has a p.f. (i.e. $P(k)$ holds)

(IS) Consider $n+1$. If $n+1$ is prime,
then $n+1 = n+1$ is a p.f. of $n+1$.

If $n+1$ is not prime then it
can be factored:

$$n+1 = a \cdot b \quad \text{where } a, b \in \{2, 3, \dots, n\}$$

neither are 1, hence neither are

by the IH, a and b have p.f.'s

$$a = p_1 p_2 \dots p_e$$

$$b = q_1 q_2 \dots q_m$$

all primes

but then $n+1 = p_1 p_2 \dots p_e q_1 q_2 \dots q_m$ is p.f.
hence $P(n+1)$ holds.

By (strong) induction, $\forall n \in \{2, 3, \dots\}$ $P(n)$ holds, i.e. every $n \geq 2$ has a p.f.

Multiple Base Cases

- Sometimes need to check more than one base case in order to make a valid IH/IS.
- esp often for recursively defined sequences

Ex: define a sequence x_n by:

$$x_1 = 2$$

$$x_2 = 3$$

$$x_n = 3x_{n-1} - 2x_{n-2} \text{ if } n \geq 3$$

Prop'n ($\forall n \in \mathbb{N}$) ($x_n = 2^{n-1} + 1$)

Pf: (BCs): if $\underline{n=1}$: $x_1 = 2 = 2^{1-1} + 1$ ✓ (98)
 if $\underline{n=2}$: $x_2 = 3 = 2^{2-1} + 1$ ✓

(IH) Fix $n \geq 2$ and assume $\forall k \in \{1, 2, \dots, n\}$

We have $x_k = 2^{k-1} + 1$

(notice: IH fixes $n \geq$ last verified BC — that's our "spring off" point)

(IS) Then: $x_{n+1} = \overset{\text{def'n of sequence}}{3x_n - 2x_{n-1}}$ (*)

$$\begin{aligned} &\stackrel{\text{IH}}{=} 3(2^{n-1} + 1) - 2(2^{n-2} + 1) \\ &= 3 \cdot 2^{n-1} + 3 - 2^{n-1} - 2 \\ &= 2 \cdot 2^{n-1} + 1 \\ &= 2^n + 1 \\ &= 2^{(n+1)-1} + 1 \quad \checkmark \end{aligned}$$

by induction, $x_n = 2^{n-1} + 1$ holds for all $n \in \mathbb{N}$

[note] — we really needed to check both $n=1$ and $n=2$ as BCs.

- if we only checked $n=1$ and
let our I.H. be: "Fix $n \geq 1$ and assume
 $\forall k \in [1, \dots, n] \quad x_k = 2^{k-1} + 1$ "

then step (*) would have been
unjustified for $n=1$, since in this case
(*) would be:

$$x_{1+1} = 3x_1 - 2x_0 \quad \swarrow \text{undefined!!}$$

- can cook up false induction proofs
that play on this wuc.

e.g. "Prop'n" let the sequence x_n
be defined as above. Then $\forall n \in \mathbb{N}$ we have:

$$x_n = 2^{n+1} - 2$$

"PF" (BC) if $n=1$, then $x_1 = 2 = 2^{1+1} - 2 \checkmark$

(IH) Fix $n \geq 1$ and assume $\forall k \in [1, \dots, n]$
that $x_k = 2^{k+1} - 2$

(IS) Then: $x_{n+1} = 3x_n - 2x_{n-1} \quad (*)$
$$= 3(2^{n+1} - 2) - 2(2^n - 2)$$

$$\begin{aligned} &= 3 \cdot 2^{n+1} - 6 - 2^{n+1} + 4 \\ &= 2 \cdot 2^{n+1} - 2 \\ &= 2^{(n+1)+1} - 2 \quad \checkmark \end{aligned}$$

By induction, the identity is "proved" $\forall n \in \mathbb{N}$.

of course, we can verify the identity is wrong even for $n=2$.

$$x_2 = 3 \neq 6 = 2^{2+1} - 2.$$

- the issue is exactly that (*) is not justified when $n=1$, but in our $\mathbb{N}$ we allow the possibility $n=1$ since we've only verified $n=1$ in our BC.

PMI, PSMI, WOP

"Theorem" (Well-ordering principle)
If $X \subseteq \mathbb{N}$ and $X \neq \emptyset$ then X has a least element (i.e. $(\exists x_0) (\forall x \in X) (x_0 \leq x)$)

- this "theorem" is intuitively obvious and is often taken as an axiom for $\mathbb{N}$.

- e.g. if $X = \mathbb{N}$ then least el't is 1
- if $X = \{2, 4, 6, \dots\}$ then least el't is 2.
- if $X = \{n \in \mathbb{N} \mid (\exists k \in \mathbb{N})(k > 5 \wedge n = k^2)\}$
 $= \{36, 49, 64, \dots\}$

then least el't is 36.

- though it's obvious, one can actually prove WOP by strong induction.

PF: We want to prove:

$$(\forall X \in P(\mathbb{N})) (X \neq \emptyset \Rightarrow X \text{ has a least el't})$$

- so fix $X \in P(\mathbb{N})$
- we argue by contrapositive
 - assume X has no least el't
 - we prove $X = \emptyset$ by strong induction
- specifically we prove:

$$(\forall n \in \mathbb{N})(n \notin X)$$

call this $P(n)$

(BC) $P(1)$ is true (i.e. $1 \notin X$) because

(f we had $1 \in X$ it would be the least element of X (1 is least in N !)

(IIt) Fix $n \in N$. Assume $\forall k \in \{1, 2, \dots, n\}$ we have $k \notin X$ (i.e. $P(k)$ holds)

(IS) Consider $n+1$. IF $n+1 \in X$ it would be least in X since by the IIt

$$1 \notin X \wedge 2 \notin X \wedge 3 \notin X \wedge \dots \wedge n \notin X$$

Hence $n+1 \notin X$, since X has no least el't.

by strong induction, $(\forall n \in N) (n \notin X)$ holds
i.e. $X = \emptyset$.

We've proved WOP. ✓

We just showed

$$PSMI \Rightarrow WOP$$

In fact, PSMI and WOP are equivalent

i.e. can also prove $WOP \Rightarrow PSMI$

and more: both are equivalent to

PMI.